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Transmitted herewith for filing is the patent application of

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For: RANDOM NUMBER GENERATION APPARATUS AND RANDOM NUMBER  
GENERATION METHOD

Enclosed are:

- ☒ Specification and Claim(s).
- ☒ Oath or Declaration (Unexecuted).
- ☒ FOUR sheet(s) of drawings.
- ☐ An assignment of the invention to \_\_\_\_\_.
- ☒ Copy of One priority application(s).
- ☐ Preliminary Amendment

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## Title of the Invention

Random Number Generation Apparatus and Random Number Generation Method

## Background of the Invention

### Field of the Invention

The present invention relates to a random number generation apparatus and a random number generation method for generating a random number sequence.

### Description of the Related Arts

As a conventional random number generation method on a computer, there can be exemplified the linear congruence method or multiplication congruence method and method using a shift register or DES (data encryption standard) which is one of the data encryption standards.

A random number sequence generated by the aforementioned methods inevitably has a regularity and its periodicity is a short. Accordingly, it is not proper to use such a random number sequence for generating an encryption key and a seed for generating an encryption key or for encryption of a message.

### Summary of the Invention

It is therefore an object of the present invention to provide a random number generating apparatus and a random number generating method for generating a random

number sequence having a long periodicity.

The random number generation apparatus according to the present invention includes: pick-up means, digital image conversion means for converting into a digital image a pick-up signal output from the pick-up means, storage means for storing the digital image as pixel values, and random number generating means for extracting a digital data from pixel values of a plurality of pixels in the digital image of the pick-up signal output when no subject is present from the pick-up means stored in the storage means and generating a random number from the digital data correlated to the plurality of pixels.

In the random number generating apparatus having the aforementioned configuration, a pick-up signal output from the pick-up means is converted into a digital image by the digital image conversion means and pixel values of the digital image are stored in the storage means. The random number generating apparatus extracts a digital data from pixel values of a plurality of pixels within the digital image of the pick-up signal output when no subject is present from the pick-up means stored in the storage means, so that the random number generation means generates a random number from the digital data correlated to the plurality of pixels.

Since there is no regularity in the pixel values of the respective pixels of the digital image obtained when no subject is present, the random number generated by the random number generating apparatus has a long periodicity.

Moreover, in the random number generating apparatus according to the present

invention, in order to solve the aforementioned problem, a pick-up signal output from the pick-up means when no subject is present is converted into a digital image and a digital data is extracted from pixel values of a plurality of pixels within the digital image, so that a random number is generated from the digital data correlated to the plurality of pixels.

Since there is no regularity in the pixel values of the respective pixels of the digital image obtained when no subject is present, the random number generated by the random number generating method has a long periodicity.

#### Brief Description of the Drawings

Fig. 1 is a block diagram showing a fingerprint identification apparatus including a random number generation apparatus according to the present invention as an encryption block.

Fig. 2 is a block diagram showing a configuration of a pick-up block of the aforementioned fingerprint identification apparatus.

Fig. 3 shows a binary image of a fingerprint obtained in the aforementioned fingerprint identification apparatus.

Fig. 4 shows an image made up by the least significant one bit of the gray scale image of the fingerprint obtained in the aforementioned fingerprint identification apparatus.

Fig. 5 is a flowchart showing a random number generation step, a prime number

generation step, and a key generation step.

Fig. 6 is a block diagram showing a binary image generator in an image processing block of the aforementioned fingerprint identification apparatus.

### Detailed Description of Preferred Embodiments

Hereinafter, a detailed explanation will be given on an embodiment of the present invention with reference to the attached drawings. As shown in Fig. 1, this embodiment is a fingerprint identification apparatus including an A/D converter 1, an encryption block 2 having a random number generator 3 and encryption means 4, a CPU 5, a memory 6, an interface block 7, and a fingerprint identifier 8. Here, the A/D converter 1, the random number generator 3, and the memory 6 constitute an example of configuration of a random number generation apparatus according to the present invention.

The fingerprint identification apparatus constitutes a personal identification apparatus for identifying a person according to a fingerprint image fetched by a pick-up block 10. In this fingerprint identification apparatus, when a desired person is identified according to a fingerprint image as a living body information, an encryption key is generated according to a random number sequence generated in the random number generator 3 and a plain text is encrypted.

The pick-up block 10 is constructed so as to pick-up a fingerprint as a living body information. More specifically, as shown in Fig. 2, the pick-up block 10 includes

a light source 11, a prism 12, and pick-up means 13.

The prism 12 has a triangular cross section. The light from the light source 11 is incident from the first face 12a, reflected from a subject placed on the second face 12b, and goes out from the third face 12c. Here, the subject is a fingerprint of a finger 100 for identifying an individual person. The pick-up means 13 is arranged at a position so as to detect the light emitted from the third face 12c. The pick-up means 13 is, for example, a CCD (Charge-Coupled Device) camera.

In the pick-up block 10 having the aforementioned configuration, when the finger 100 is placed on the second face 12b of the prism 12, the light emitted from the light source 11 comes into the prism through the first face 12a and is reflected irregularly by the convex portion of the fingerprint of the finger 100 on the second face 12b or reflected totally by the concave portion. These reflected lights go out from the third face 12c and form an image in the pick-up means 13. Thus, in the pick-up means 13, the convex portion of the finger 100 is picked up as a dark portion and the concave portion is picked up as a bright portion. The pick-up means 13 outputs a pick-up signal as a pick-up information.

The pick-up signal output from the pick-up block 10 is sampled at an appropriate time interval and converted by the A/D converter 1 into a digital image of a size, for example,  $256 \times 128$ . In this embodiment, the A/D converter 1 performs an 8-bit conversion. Thus, pixel values of the pixels constituting an image are digital data expressed by 256 gradations from 0 to 255. The digital image obtained in this A/D

converter 1 is stored in the memory 6. Hereinafter, a digital image whose pixel value is expressed by multiple bits such as 8 bits will be referred to as a gray scale image."

According to the gray scale image, the image processor 20 generates a binary image. For example, the image processor 20 fetches the gray scale image at an appropriate timing and using an appropriate binarization method, generates a binary image in which pixel value of each 8-bit pixel has been converted into '0' or '1'. The binarization method may be a comparison between an average of pixel values of the entire image and pixel values of the respective pixels or a moving average method, i.e., comparison between a pixel value of a pixel to be considered and an average of pixel values of a plurality of pixels located in a predetermined range from the considered pixel. For example, the fingerprint image picked up in the pick-up block 10 is made into a binary image as shown in Fig. 3 by the moving average method. In Fig. 3, the black portions represent convex portions of the fingerprint and the white portions represent concave portions of the fingerprint.

The binary image thus generated is subjected to a pre-processing such as a thinning and then processes such as registration and identification are performed. It should be noted that the binary image generation from the gray scale image by the aforementioned moving average method will be detailed later.

The fingerprint identifier 8 identifies the binary image. For example, the fingerprint identifier 8 identifies a registered image information on the fingerprint information which has been fetched in advance with the binary image of the fingerprint

picked up by the pick-up block 10. According to the identification result in the fingerprint identifier 8, the fingerprint identification apparatus identifies a desired individual.

It should be noted that the CPU 5 is control means for controlling respective components constituting the fingerprint identification apparatus.

As has been described above, the fingerprint identification apparatus identifies a fingerprint from a digital image picked up by the pick-up block 10 to identify a desired individual. When an individual is identified by such a fingerprint identification process, the fingerprint identification apparatus encrypts a plain text using a private key. This encryption using a private key is performed according to a prime number obtained according to a random number sequence generated by the random number generator 2.

Next, explanation will be given on the process how the encryption block 2 causes the random number generator 3 to generate a random number sequence and the encryption means 4 to perform encryption using an encryption key according to the random number sequence. It should be noted that although the random number generator 3 is constructed to generate a random number sequence from the aforementioned gray scale image or the binary image, explanation will be given on a case of generating a random number sequence according to a gray scale image.

In the pick-up block 10, when an image is taken in without placing a finger on the prism 12, a pick-up signal output from the pick-up means 13 is overlapped with a

noise. As a result, the least significant bit of the gray scale image obtained by digital conversion in the A/D converter 1 shows a value of irregular '0' or '1'. For example, similar irregular values are shown for a binary image. Accordingly, in the gray scale image, it is possible to obtain a random number sequence consisting of '0' and '1' and having an arbitrary length from a bit sequence of an arbitrary length starting at an appropriate position as a start address. For example, in the gray scale image, when it is assumed that the least significant bit '0' represents black and the least significant bit '1' represents white, it is possible to obtain a binary image as shown in Fig. 4. As shown in this Fig. 4, the least significant bits of the gray scale image have no regularity.

According to the random number sequence obtained by the random number generator 3, the encryption block 2 generates an encryption key or seed as an origin of the encryption key and performs encryption in the encryption means 4.

given on this case.

Moreover, the RSA encryption method creates a 384-bit, 512-bit, or 1024-bit key for encryption. Here, explanation will be given on a case using the 512-bit key. The outline of the RSA encryption method is as follows.

In the RSA encryption method, from two prime number  $p$  and  $q$  and one of the public keys  $E$  (public exponent), using Equations (1) and (2), the other public key, i.e., the public key  $N$  (modulus) and a private key  $D$  (private exponent) will be obtained.

$$N = p \times q \dots (1)$$

$$D = E^{-1} \bmod \{(p - 1) \times (q - 1)\} \dots (2)$$

Here, the public key  $E$  and the multiple of  $(p - 1)$  and  $(q - 1)$  are mutually prime. If a message (plain text) is  $M$  and an encrypted message is  $C$ , then relations expressed by Equations (3) and (4) are satisfied.

$$C = M^E \bmod N \dots (3)$$

$$M = C^D \bmod N \dots (4)$$

The public key  $N$  is a very large 512-bit number and it is very difficult to factorize it into prime factors and accordingly, the addressee cannot obtain the previous message  $M$  from the encrypted message  $C$  unless the addressee knows the private key  $D$ . Moreover, in order to add a digital signature to the message  $C$  when sent to the addressee, the addresser encrypts the message  $C$  having his/her signature using his/her private key  $D$  according to Equation (4) when sending the message  $M$ . The addressee decodes the message using the public key  $E$  and the public key  $N$  of the

addresser according to Equation (3) and confirms that the signature of the addresser is added.

This is the outline of the RSA encryption method. In the encryption means 4 employing the RSA encryption method, a 512-bit key is required. A random number sequence generated in the random number generator 3 is used for creating such a 512-bit key. Such a 512-bit key can be generated by generation of a random number sequence as follows.

Since the key length is 512 bits, firstly, the random number generator 3 generates two 256-bit random number. These two random numbers serve as seeds, i.e., initial values for finding two prime numbers.

As has been described above, when generating a random number, the fingerprint identification apparatus takes in an image without placing a finger 100 on the prism 12 and obtains a gray scale image as a digital image formed by the A/D converter 1. The fingerprint identification apparatus stores the gray scale image on memory 6 as having size of 256 pixels in the horizontal direction and 128 pixels in the vertical direction in which each pixel value is expressed by 8 bits. It should be noted that simultaneously with such a gray scale image, the fingerprint identification apparatus fetches a binary image from this gray scale image by the image processor 20. The fingerprint identification apparatus stores the binary image on memory 6 as having a size of 256 pixels in the horizontal direction and 128 pixels in the vertical direction in which each pixel is expressed by 1 bit.

As has been described above, the least significant bits of the pixel value of pixels in the gray scale image have no regularity. Accordingly, by extracting the least significant bits of the pixel values for a plurality of pixels, it is possible to generate a random number having a long periodicity. The random number generator 3 generates a random number by extracting the least significant bits of pixel values of a plurality of pixels constituting a predetermined area starting at a pixel located at a start address. Here, the start address is an information indicating a position of a pixel where the least significant bit extraction is started.

More specifically, pixels are scanned in the horizontal direction starting at the start address so as to extract the least significant bit value of the pixels, i.e., '0' or '1'. Assuming  $i$  for the horizontal direction address and  $j$  for the vertical direction address, an arbitrary point on the gray scale image is defined as  $g(i, j)$ .

For example, the start address is defined as (128, 0) and the 512 pixels are scanned from the pixel  $g(128, 0)$  to the pixel  $g(129, 255)$  and the least significant bit values are extracted to generate two 256-bit random numbers.

Moreover, it is also possible to generate a random number by defining a start address at an appropriate position instead of a predetermined position. In this case, values from 0 to 127 are expressed by 7 bits. Accordingly, by defining the start address by the horizontal address  $i$  and the vertical address  $j$  specified by the 8 bits of pixel values of the pixel  $g(0, 0)$  and the least significant 7 bits of the pixel values of the pixel  $g(0, 1)$ , values of the least significant bits of the pixel values of pixels are

extracted to generate a random number. For example, when the value expressed by an 8-bit pixel value of the pixel  $g(0, 0)$  is 100 and the value expressed by the least significant 7 bits of the pixel value of the pixel  $g(0, 1)$  is 23, the least significant bit of the pixel value of the pixel  $g(100, 23)$  is extracted to generate a random number.

Furthermore, when there is a correlation between two adjacent pixels in the horizontal direction, a particular pattern (random number) is easily generated. Taking this into consideration, the least significant bit of the pixel value can be extracted. For example, scan is performed in the vertical direction and the least significant bit of the pixel value is extracted. Moreover, it is possible to perform an exclusive OR operation between two adjacent pixels in the vertical direction to extract a 1-bit data. Alternatively, it is possible to perform an image take-in twice and perform an exclusive OR operation between two images so as to extract a 1-bit data.

As has been described above, the random number generator 3 generates a complete random number having a long periodicity by extracting the least significant bit of pixel values of pixels. The encryption means 4 generates two prime number  $p$  and  $q$  from the two random numbers generated by the random number generator 3. As shown in fig. 5, the encryption means 4 generates an encryption key through a prime number generation process and a key generation process.

Firstly, as shown in Fig. 5, according to a random number generated by the least significant bits of pixel values (gray scale data) of a gray scale image in step 1, the encryption means 4 generates a prime number in the prime number generation

process of steps S2 to S5. It should be noted that the process described below is performed for each of the two random numbers  $p$  and  $q$ .

As shown in step S2, the encryption means 4 sets the most significant bit and the least significant bit to '1'. Thus, the random number generated in step S1 has a length of 256 bits and is an odd number.

Next, in step S3, the encryption means 4 performs division of the random number using all the prime numbers smaller than 256 to determine whether the random number can be divided by all the prime numbers without a remainder. Here, unless the random number can be divided by all the prime numbers without a remainder, the encryption means 4 passes control to step S4, and if the random number can be divided by all the prime numbers without a remainder, the encryption means 4 passes control to step S5.

In step S4, the encryption means 4 uses the Rabin-Miller method which is a representative probability prime number checking method so as to further check whether the random number which has been subjected to division tests by all the prime numbers smaller than 256 in step S3 is a prime number. Here, if the number is determined to be a prime number, control is passed to step S6, and otherwise, control is passed to step S5.

In step 5 which is performed even if the random number is divided by prime numbers without a remainder in step S3, the encryption means 4 subtracts 2 from the value of the random number  $p$  (or the random number  $q$ ). Then, control is passed to

step S3, where the encryption means 4 again checks whether the random number subtracted by 2 can be divided by all the prime numbers smaller than 256 without a remained so as to perform the aforementioned processes of step S3 or step S5 and after.

In step S6, as a key generation step, the encryption means 4 fetches a public key  $N$  from the aforementioned Equation (1) according to the two prime numbers  $p$  and  $q$ , and from this public key  $N$  and a public key  $E$  appropriately selected, obtains a private key  $D$  satisfying the aforementioned Equation (2). For example, the encryption means 4 obtains the private key  $D$  satisfying the aforementioned Equation (2) by an extended Euclidean algorithm.

As has been described above, in the random number generated by the random number generation step, the most significant bit and the least significant bit are set to '1' in the prime number generation step and the key generation step. Thus, the random number has a 256-bit length and is an odd number. This random number is successively divided by all the prime numbers smaller than 256 and it is confirmed that the random number cannot be divided without a remainder by any of the prime numbers. The random number which has been confirmed that it cannot be divided by any of the prime numbers smaller than 256 is then subjected to a check using the Rabin-Miller method which is a representative probabilistic primality test to determine whether the random number tested is a prime number. Here, if the number is determined not to be a prime number, the random number tested is subtracted by 2 and

then again subjected to a check to determine whether the number is a prime number. If the random number is determined to be a prime number, the random number is used to obtain the private key  $D$  satisfying the Equation (2) from the public key  $N$  calculated from the Equation (1) and the public key  $E$ .

As has been described above, the encryption block 2 causes the random number generator 3 to generate a complete random number having a long periodicity and the encryption means 4 to generate a prime number according to this random number, so that the prime number is used to generate the private key D as an encryption key. The fingerprint identification apparatus has private key custody means for keeping the private key D in custody. The private key D thus generated is stored, for example, in the memory 6 functioning as the private key custody means and thus kept in custody within the fingerprint identification apparatus.

The encryption block 2 uses the private key D to encrypt a message (plain text). The message is added by a digital signature as follows in the encryption block 2.

The fingerprint identification apparatus identifies a binary image obtained when a finger is placed on the prism 12 in the fingerprint identification block 8 and identifies the individual. When the individual is identified, the encryption block 2 uses the private key D to encrypt the message. Here, the fingerprint identification apparatus is connected via the interface block 6 to a personal computer (not depicted) and the message has been transmitted via the interface block 6 from the personal computer.

The fingerprint identification apparatus adds a digital signature to the encrypted

message in the encryption block 2 and sends the message back to the personal computer.

The personal computer transmits to a desired addressee the encrypted message having the digital signature via a network.

As has been described above, the fingerprint identification apparatus, upon identification of a desired individual, uses an encryption key to encrypt a message and sends the encrypted message to a desired addressee.

As has been described above, this fingerprint identification apparatus uses the least significant bits of a gray scale image obtained in the pick-up block 10 when no finger 100 is placed on the prism 12 and obtains a random number having a long periodicity. According to such a random number, the fingerprint identification apparatus generates a prime number to be used in encryption, thus providing an encryption with a high reliability.

Furthermore, the fingerprint identification apparatus stores the private key D used for encryption, in custody means dedicated for a private key and performed encryption without showing the private key D to an external apparatus such as a personal computer connected. Thus, it is possible to provide an encryption with a high reliability. That is, an encryption is performed entirely within the fingerprint identification apparatus while keeping the private key D in the fingerprint identification apparatus, so that the private key D will not be read by a third party and the sequence of processes for random number generation and encryption can be

performed within one and the same fingerprint identification apparatus. Thus, this encryption has an improved security.

It should be noted that in the aforementioned embodiment, an explanation has been given on a case of generating a random number from the least significant bits of pixel values of a gray scale image. However, the fingerprint identification apparatus can also generate a random number according to pixel values of the respective pixels of a binary image, and can generate a random number according to pixel values of respective pixels of a binary image as follows. Here, it is assumed that the horizontal direction address is  $i$  and the vertical direction address is  $j$ , and an arbitrary pixel on the binary image is  $b(i, j)$ .

For example, similarly as in the aforementioned gray scale image, when the start address is  $(128, 0)$ , the random number generator 3 extracts pixel values of respective pixels from pixel  $b(128, 0)$  to pixel  $b(129, 255)$  and generate two 256-bit random numbers.

Moreover, the random number generator 3 can generate a random number at an arbitrary start address instead of a predetermined position on the screen. For example, the random number generator 3 uses as a start address the horizontal address  $i$  and the vertical address  $j$  specified by the pixel values of pixels from pixel  $b(0, 0)$  to pixel  $b(0, 7)$  and pixel values of respective pixels from pixel  $b(0, 8)$  to pixel  $b(0, 14)$  and extracts pixel values of pixels to generate a random number. For example, when the value specified by the pixel values of pixels from pixel  $b(0, 0)$  to pixel  $b(0, 7)$  is 100 and the

value specified by the pixel values of pixels from pixel b(0, 8) to pixel b(0, 14) is 23, the random number generator 3 extracts pixel values starting at pixel b(100, 23) so as to generate a random number.

Moreover, similarly as in the gray scale image, it is possible to extract pixel values by scanning in the vertical direction, to extract a one-bit data by the exclusive OR operations between two pixels adjacent in the vertical direction, and to perform take-in of an image twice and perform the exclusive OR operation between the two images so as to extract a one-bit data. By such extracts, the random number generator 3 can generate a more complete random number.

According to the two random numbers generated according to the binary image in the random number generator 3, the encryption means 4 generates an encryption key by the prime number generation process and the key generation process shown in Fig. 5. That is, according to the random number based on the pixel values (binary data) of the binary image generated in step S1, an encryption key is generated through the prime number generation process and the key generation process in the steps S2 to S6.

It should be noted that as shown in Fig. 6, the image processor 20 includes a binary image generation block for generating a binary image from a gray scale image. This image processor 20 is constructed corresponding to the moving average method. In this embodiment, explanation will be given on binarization performed using an average value of 7 pixels in the vertical direction and 7 pixels in the horizontal direction around a center pixel ( $7 \times 7$  pixels).



direction is calculated. Outputs from the first to the seventh horizontal direction summing blocks 28, 29, 30, 31, 32, 33, and 34 are added in the adder 35 constituting the vertical direction summing block and then input to the divider 36.

The divider 36 divides the total by the number 49 of the pixels added in the horizontal direction and the vertical direction so as to calculate a binary threshold value. The calculated value is compared by the comparator 37 to a binary threshold value of the fourth horizontal direction summing block 31 for binarization.

By the aforementioned configuration, the binary image generation block generates a binary image from the gray scale image.

The random number generator 3 can generate a random number as has been described above according to pixel values of respective pixels of the binary image thus generated by the binary image generation block.

Moreover, the fingerprint identifier 8 identifies a fingerprint according to the binary image generated by the binary image generation block.

The random number generation apparatus according to the present invention includes: digital image conversion means for converting a pick-up signal output from pick-up means, into a digital image; storage means for storing the digital image as pixel values; and random number generation means for extracting a digital data from pixel values of a plurality of pixels in a digital image of a pick-up signal output, when no subject is present, from pick-up means stored in the storage means and generating a random number from the digital data correlated to a plurality of pixels. The pick-up



function of encrypting a plain text employs the random number generation method for generating such a random number, so as to generate an encryption key within the apparatus and keep the encryption key generated, in custody within the apparatus. Thus, it is possible to perform encryption with an improved safety.

What is claimed is:

1. A random number generation apparatus comprising:  
pick-up means,  
digital image conversion means for converting into a digital image a pick-up  
signal output from said pick-up means,  
storage means for storing said digital image as pixel values, and  
random number generating means for extracting a digital data from pixel values  
of a plurality of pixels in said digital image of said pick-up signal output when no  
subject is present from said pick-up means stored in said storage means and generating  
a random number from said digital data correlated to said plurality of pixels.

2. A random number generating apparatus as claimed in Claim 1, wherein  
said pixel values are expressed by 2 or more bits, and  
said random number generating means extracts as said digital data binary data  
of the least significant bits of a bit string expressing said pixel values and generates a  
random number from said binary data correlated to said plurality of pixels.

3. A random number generating apparatus as claimed in Claim 1, wherein  
said digital image is a binary image in which said pixel values are expressed as 1-bit  
binary data, and  
said random number generating means extracts said binary data as said digital



8. A random number generating apparatus as claimed in Claim 6, wherein said apparatus further comprises custody means for keeping said encryption key in custody.

9. A random number generating apparatus as claimed in Claim 6, wherein said apparatus has said encryption block and is provided in an individual identification apparatus according to said digital image correlated to a subject picked up by said pick-up means stored in said storage means,

said individual identification apparatus performing an encryption of a plain text using said encryption key by said encryption block when a desired individual is identified.

10. A random number generating apparatus as claimed in Claim 9, wherein said subject is a fingerprint constituting a living body information.

11. A random number generating method comprising steps of  
converting into a digital image a pick-up signal output from pick-up means  
when no subject is present,

extracting a digital data from pixel values of a plurality of pixels in said digital image, and

generating a random number from the digital data correlated to said plurality of pixels.

12. A random number generating method as claimed in Claim 11, wherein said pixel values are expressed by two or more bits, and

a binary data of the least significant bit of a bit string expressing said pixel values is extracted as said digital data, so that a random number is generated from said binary data correlated to said plurality of pixels.

13. A random number generating method as claimed in Claim 11, wherein said digital image is a binary image in which said pixel values are expressed as 1-bit binary data, and

said binary data is extracted as said digital data, so that a random number is generated from said binary data correlated to said plurality of pixels.

14. A random number generating method as claimed in Claim 11, wherein said digital data is extracted from pixel values of a plurality of pixels constituting a predetermined area within said digital image.

15. A random number generating method as claimed in Claim 11, wherein said digital data is extracted from pixel values of said plurality of pixels at an arbitrary

position within said digital image.

16. A random number generating method as claimed in Claim 11, said method being used in an encryption method for generating an encryption key according to said random number or according to a data obtained from said random number.

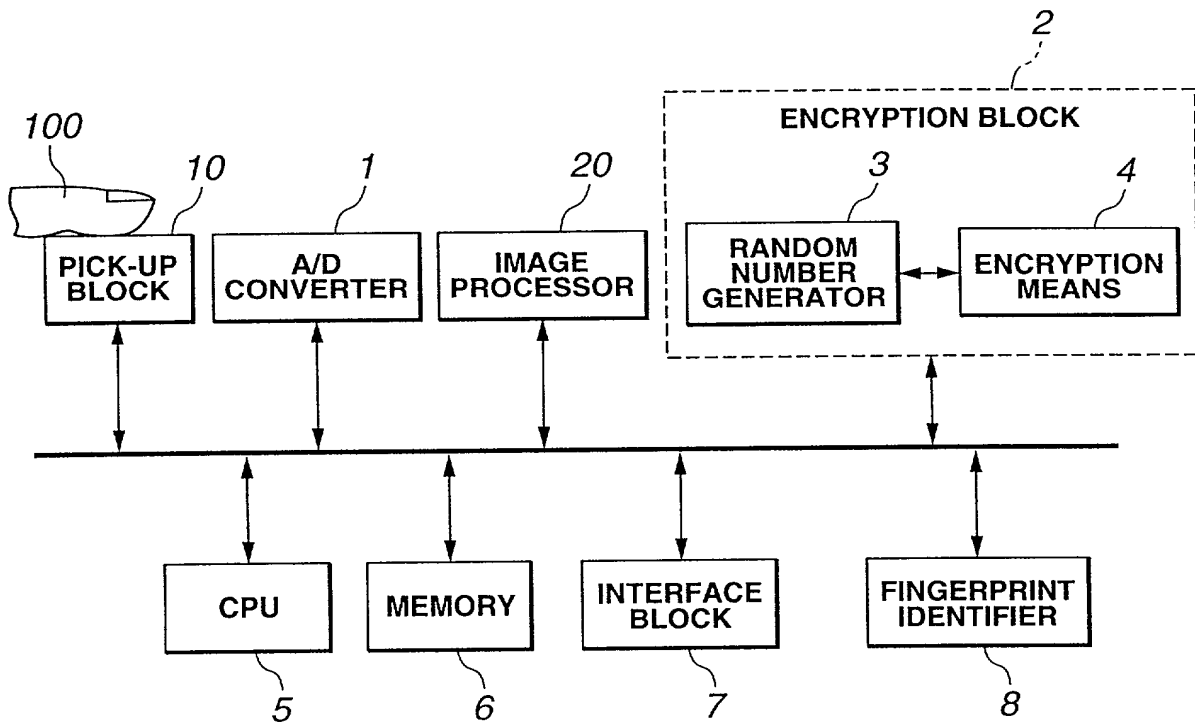
17. A random number generating method as claimed in Claim 16, wherein said encryption method is the RSA encryption method for generating said encryption key and said two prime numbers are generated according to said random number so that said two prime numbers are used for generating said encryption key.

18. A random number generating method as claimed in Claim 16, wherein in said encryption method, said encryption key is kept in custody within the apparatus which performs the encryption according to said encryption method.

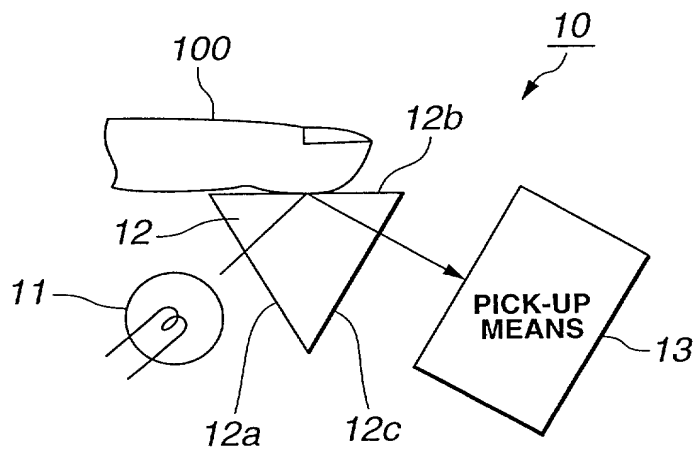
19. A random number generating method as claimed in Claim 16, wherein in said encryption method, a plain text is encrypted by said encryption key when a desired individual is identified according to said digital image correlated to a subject picked up by said pick-up means.

20. A random number generating method as claimed in Claim 19, wherein said subject is a fingerprint constituting a living body information.

[illegible]



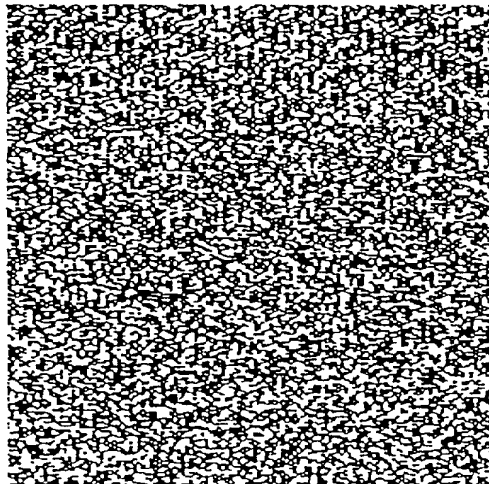
**FIG.1**



**FIG.2**



**FIG.3**



**FIG.4**

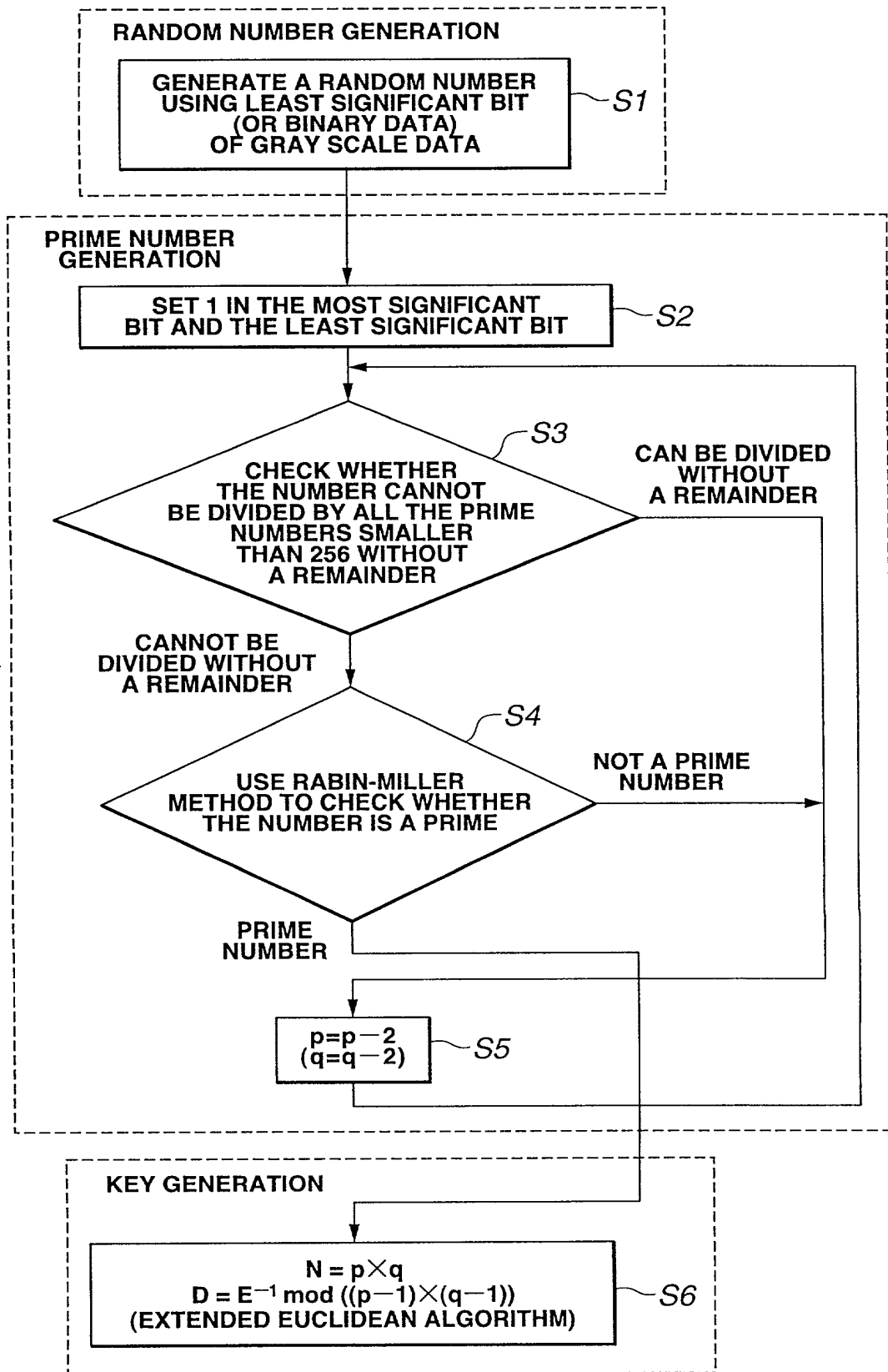


FIG.5

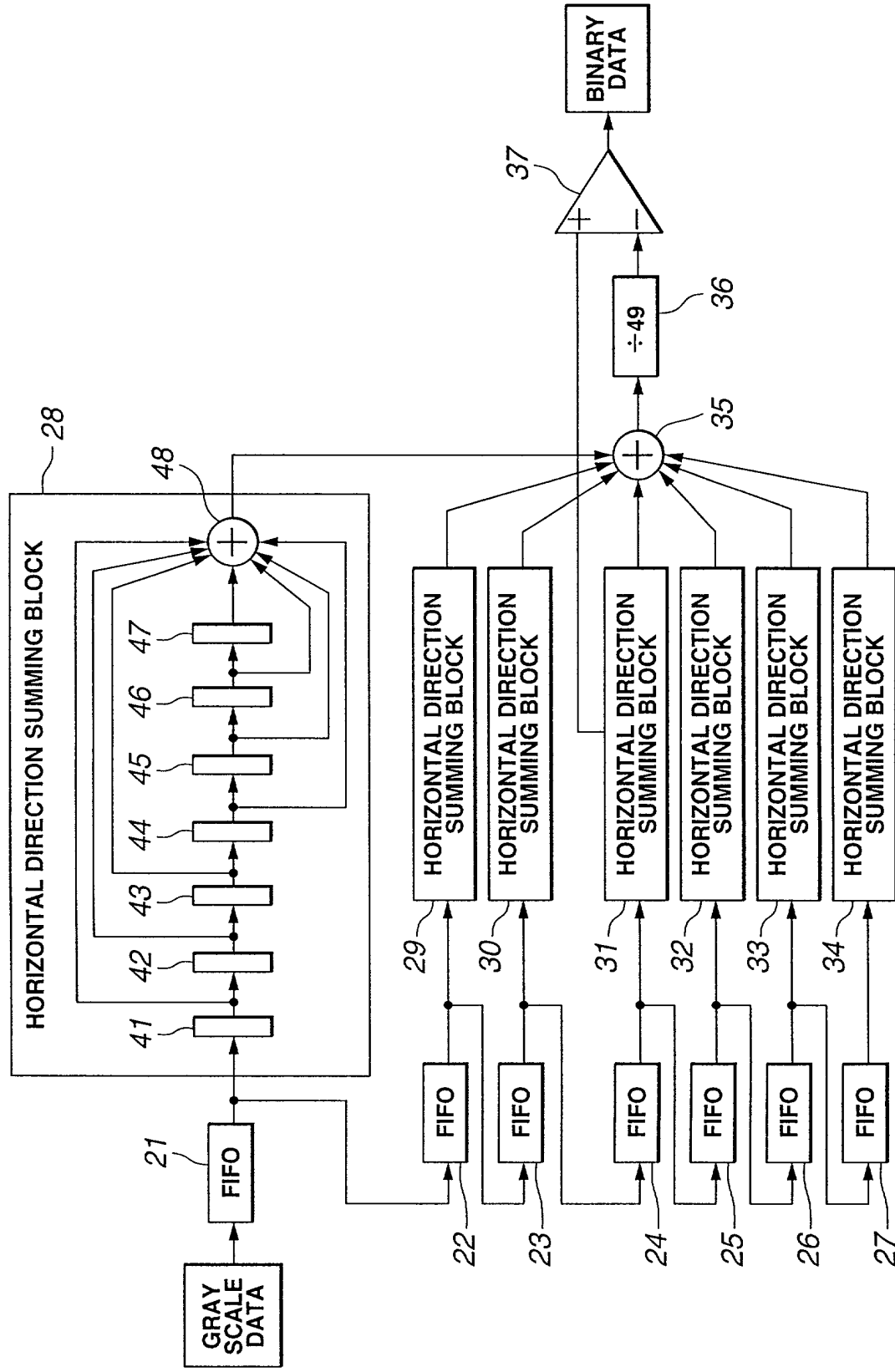


FIG. 6

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<b>DECLARATION FOR UTILITY OR DESIGN PATENT APPLICATION</b>		<b>Attorney Docket No.</b>		80001-1887	
		<b>First Named Inventor</b>		Takuya WADA	
		<b>COMPLETE IF KNOWN</b>			
		<b>Application No.</b>			
		<b>Filing Date</b>			
<input type="checkbox"/> Declaration Submitted with or Initial filing		<input checked="" type="checkbox"/> Declaration submitted after initial filing		<b>Group Art Unit</b>	
				<b>Examiner Name</b>	

As a below named inventor, I hereby declare that:

My residence, post office address, and citizenship are as stated below next to my name.

I believe I am the original, first and sole inventor (only if one name is listed below) or an original, first and joint inventor (if plural names are listed below) of the subject matter which is claimed and for which a patent is sought on the invention entitled:

**RANDOM NUMBER GERATION APPARATUS AND RANDOM NUMBER GENERATION METHOD**  
*(Title of the Invention)*

the specification of which

☒ is attached hereto

or

☐ was filed on \_\_\_\_\_, as United States Application Number or PCT International Application Number: \_\_\_\_\_ and was amended on \_\_\_\_\_ (if applicable).

I hereby state that I have reviewed and understand the contents of the above identified specification, including the claims, as amended by any amendment specifically referred to above.

I acknowledge the duty to disclose information which is material to patentability as defined in Title 37, Code of Federal Regulations § 1.56.

I hereby claim foreign priority benefits under Title 35, United States Code § 119 (a)-(d) of any foreign application(s) for patent or inventor's certificate, or § 365(a) of any PCT international application which designated at least one country other than the United States of America, listed below and have also identified below, by checking the box, any foreign application for patent or inventor's certificate, or of any PCT international application having a filing date before that of the application on which priority is claimed.

Prior Foreign Application Number(s)	Country	Foreign Filing Date (MM/DD/YY)	Priority Not Claimed	Certified Copy Attached	
				YES	NO
P11-226555	Japan	08/10/99		<input checked="" type="checkbox"/>	<input type="checkbox"/>
				<input type="checkbox"/>	<input type="checkbox"/>
				<input type="checkbox"/>	<input type="checkbox"/>

☐ Additional foreign application numbers are listed on a supplemental priority data sheet PTO/SB/02B attached hereto:

I hereby claim the benefit under Title 35, United States Code § 119(e) of any United States provisional application(s) listed below.

Application Number (s)	Filing Date (MM/DD/YY)	<input type="checkbox"/> Additional provisional application numbers are listed on a supplemental priority data sheet PTO/SB/02B attached hereto.

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**DECLARATION - Utility Or Design Patent Application**

I hereby claim the benefit under Title 35, United States Code §120 of any United States application(s) of any PCT International application designating the United States of America, listed below and, insofar as the subject matter of each of the claims of this application is not disclosed in the prior United States or PCT International application in the matter provided by the first paragraph of Title 35, United States Code §112, I acknowledge the duty to disclose information which is material to patentability as defined in Title 37, Code of Federal Regulations § 1.56 which became available between the filing date of the prior application and the national or PCT international filing date of this application.

U.S. Parent Application Number	PCT Parent Number	Parent Filing Date (MM/DD/YYYY)	Parent Patent Number (if applicable)

☐ Additional U.S. or PCT international application numbers are listed on a supplemental priority data sheet PTO/SB/02B attached hereto

As a named inventor, I hereby appoint the following registered practitioner(s) to prosecute this application and to transact all business in the Patent and Trademark Office connected therewith: ☐ Customer Number

or

☒ Registered practitioner(s) name/registration number listed below

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Glenn E. Forbis	40,610	Monica E. Millner	42,894
Ronald P. Kananen	24,104	David K. Benson	42,314
Matthew J. Russo	41,282	Christopher M. Tanner	41,518
G. Thomas Williams	42,228	Robert S. Green	41,800
Shmuel Livnat	33,949		

☐ Additional registered practitioner(s) named on supplemental Registered Practitioner Information sheet PTO/SB/02C attached hereto.

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I hereby declare that all statements made herein of my own knowledge are true and that all statements made on information and belief are believed to be true; and further that these statements were made with the knowledge that willful false statements and the like so made are punishable by fine or imprisonment, or both, under Section 1001 of Title 18 of the United States Code and that such willful false statements may jeopardize the validity of the application or any patent issued thereon.

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Takuya				WADA			
<b>Inventor's Signature</b>						<b>Dated</b>	
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<b>Name of Second Inventor</b>		<input type="checkbox"/> A petition has been filed for this unsigned inventor					
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Yoshihiro				TSUKAMURA			
<b>Inventor's Signature</b>						<b>Dated</b>	
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